Memory 3: Caching and TLBs (Con’t), Demand Paging

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How is the Translation Accomplished?

- The MMU must attempt to translate virtual address to physical address on:
  - Every instruction fetch, Every load, Every store
  - Generate a “Page Fault” (Trap) if it encounters invalid PTE
    » Fault handler will decide what to do (more on this next lecture)
- What does the MMU need to do to translate an address?
  - 1-level Page Table
    » Read PTE from memory, check valid, merge address
    » Set “accessed” bit in PTE, Set “dirty bit” on write
  - 2-level Page Table
    » Read and check first level
    » Read, check, and update PTE at second level
  - N-level Page Table …
- MMU does page table Tree Traversal to translate each address
  - Turns a potentially fast memory access into a slow multi-access table traversal…
  - Need CACHING!

Recall: The two-level page table

- Tree of Page Tables
  - “Magic” 10b-10b-12b pattern!
- Tables fixed size (1024 entries)
  - On context-switch: save single PageTablePtr register (i.e. CR3)
- Valid bits on Page Table Entries
  - Don’t need every 2nd-level table
  - Even when exist, 2nd-level tables can reside on disk if not in use

Recall: CS61c Caching Concept

- Cache: a repository for copies that can be accessed more quickly than the original
  - Make frequent case fast and infrequent case less dominant
- Caching underlies many techniques used today to make computers fast
  - Can cache: memory locations, address translations, pages, file blocks, file names, network routes, etc…
- Only good if:
  - Frequent case frequent enough and
  - Infrequent case not too expensive
- Many important OS concepts boil down to caching! We cache:
  - Pages, Files, Virtual Memory Translations, IP Addresses…
Recall: In Machine Structures (eg. 61C) …

- Hardware Caching is the key to memory system performance for CPUs:
  - Average Memory Access Time (AMAT) = (Hit Rate x HitTime) + (Miss Rate x MissTime)
  - Where:
    - HitRate + MissRate = 1
    - MissTime = HitTime + MissPenalty
  - Examples:
    - HitRate = 90% => AMAT = (0.9 x 1) + (0.1 x 101)=11 ns
    - HitRate = 99% => AMAT = (0.99 x 1) + (0.01 x 101)=2.01 ns

Another Major Reason to Deal with Caching

- Cannot afford to translate on every access
  - At least three DRAM accesses per actual DRAM access
  - Or: perhaps I/O if page table partially on disk!
- Even worse: What if we are using caching to make memory access faster than DRAM access?
  - Solution? Cache translations!
    - Translation Cache: TLB ("Translation Lookaside Buffer")

Why Does Caching Help? Locality!

- Temporal Locality (Locality in Time):
  - Keep recently accessed data items closer to processor
- Spatial Locality (Locality in Space):
  - Move contiguous blocks to the upper levels

Recall: Memory Hierarchy

- Caching: Take advantage of the principle of locality to:
  - Present the illusion of having as much memory as in the cheapest technology
  - Provide average speed similar to that offered by the fastest technology
Recall 61C: Dealing with Hierarchy

- Used to compute access time probabilistically:
  \[
  \text{AMAT} = \text{Hit Rate}_{L1} \times \text{Hit Time}_{L1} + \text{Miss Rate}_{L1} \times \text{Miss Time}_{L1}
  \]
  \[
  \text{Hit Time}_{L1} = \text{Time to get value from L1 cache.}
  \]
  \[
  \text{Miss Time}_{L1} = \text{Time to get value from lower level (DRAM)}
  \]
  So, \[
  \text{AMAT} = \text{Hit Time}_{L1} + \text{Miss Rate}_{L1} \times \text{Miss Penalty}_{L1}
  \]
  \[
  \text{Miss Penalty}_{L1} = \text{Average Time to fetch from below L2 (DRAM)}
  \]
  
- And so on ... (can do this recursively for more levels!)

How is a Block found in a Cache?

- Block is minimum quantum of caching
  - Data select field used to select data within block
  - Many caching applications don’t have data select field
- Index Used to Lookup Candidates in Cache
  - Index identifies the set
- Tag used to identify actual copy
  - If no candidates match, then declare cache miss

Review: Direct Mapped Cache

- Direct Mapped \(2^N\) byte cache:
  - The uppermost (32 - N) bits are always the Cache Tag
  - The lowest M bits are the Byte Select (Block Size = \(2^M\))
- Example: 1 KB Direct Mapped Cache with 32 B Blocks
  - Index chooses potential block
  - Tag checked to verify block
  - Byte select chooses byte within block

Review: Set Associative Cache

- N-way set associative: N entries per Cache Index
  - N direct mapped caches operates in parallel
- Example: Two-way set associative cache
  - Cache Index selects a “set” from the cache
  - Two tags in the set are compared to input in parallel
  - Data is selected based on the tag result
**Review: Fully Associative Cache**

- **Fully Associative:** Every block can hold any line
  - Address does not include a cache index
  - Compare Cache Tags of all Cache Entries in Parallel
- **Example:** Block Size=32B blocks
  - We need N 27-bit comparators
  - Still have byte select to choose from within block

```
Cache Tag (27 bits long)  |  Byte Select
-------------------------------
|                             |
-------------------------------
```

Ex: 0x01

```
Valid Bit
```

```
Cache Tag
```

```
Byte 0  Byte 1  Byte 31  Byte 32  Byte 33  Byte 63
```

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**Administrivia (1/2)**

- Happy π Day!!!
  - 40 digits are sufficient to calculate circumference of visible universe to atomic dimensions:
  

  - Here are 40 decimal places:
    3.1415926535897932384626433832795028841971

  - Best formula for PI is from Ramanujan:

    \[
    \frac{1}{\pi} = \frac{2\sqrt{2}}{9801} \sum_{k=0}^{\infty} \frac{(4k)(1103+26390k)}{(k!)^4 396^{4k}}
    \]

  - Google announced back in 2019 (3/14/19) that Emma Haruka Iwao had just calculated π to 31,415,926,535,897 digits (new record…)

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**Administrivia (2/2)**

- Midterm 2: TOMORROW!
  - 8pm-10pm, 150 Wheeler Hall
  - You are responsible material up to and including today’s lecture (specifically, caching and basic idea of TLBs).
  - Two sheets of notes: handwritten, double-sided
- Busy working on Project 2 and Homework 4!
- Make sure to fill out the survey!
  - We want to know how we are doing
  - Also, get to consider topics for optional lecture at end of the term…

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**Where does a Block Get Placed in a Cache?**

- **Example:** Block 12 placed in 8 block cache

```
32-Block Address Space:
```

```
Block no. 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
```

**Direct mapped:**
- Block 12 can go only into block 4 (12 mod 8)

**Set associative:**
- Block 12 can go anywhere in set 0 (12 mod 4)

**Fully associative:**
- Block 12 can go anywhere
**Which block should be replaced on a miss?**

- Easy for Direct Mapped: Only one possibility
- Set Associative or Fully Associative:
  - Random
  - LRU (Least Recently Used)

**Miss rates for a workload:**

<table>
<thead>
<tr>
<th>Size</th>
<th>2-way LRU</th>
<th>Random LRU</th>
<th>4-way LRU</th>
<th>Random LRU</th>
<th>8-way LRU</th>
<th>Random LRU</th>
</tr>
</thead>
<tbody>
<tr>
<td>16 KB</td>
<td>5.2%</td>
<td>5.7%</td>
<td>4.7%</td>
<td>5.3%</td>
<td>4.4%</td>
<td>5.0%</td>
</tr>
<tr>
<td>64 KB</td>
<td>1.9%</td>
<td>2.0%</td>
<td>1.5%</td>
<td>1.7%</td>
<td>1.4%</td>
<td>1.5%</td>
</tr>
<tr>
<td>256 KB</td>
<td>1.15%</td>
<td>1.17%</td>
<td>1.13%</td>
<td>1.13%</td>
<td>1.12%</td>
<td>1.12%</td>
</tr>
</tbody>
</table>

**Review: What happens on a write?**

- **Write through:** The information is written to both the block in the cache and to the block in the lower-level memory
  - Solution: Process write through if buffer not full
  - Problem: How full is buffer?
  - Solution: Buffer size

- **Write back:** The information is written only to the block in the cache
  - Modified cache block is written to main memory only when it is replaced
  - Question is block clean or dirty?
  - WT:
    - PRO: read misses cannot result in writes
    - CON: Processor held up on writes unless writes buffered
  - WB:
    - PRO: repeated writes not sent to DRAM
    - CON: More complex

**How do we make Address Translation Fast?**

- Cache results of recent translations
  - Different from a traditional cache
  - Cache Page Table Entries using Virtual Page # as the key

**A Summary on Sources of Cache Misses**

- **Compulsory** (cold start or process migration, first reference): first access to a block
  - "Cold" fact of life: not a whole lot you can do about it unless you prefetch
  - Solution: Prefetch values before use
  - Note: If you run "billions" of instruction, Compulsory Misses are insignificant

- **Capacity:**
  - Cache cannot contain all blocks access by the program
  - Solution 1: increase cache size
  - Solution 2: change associativity

- **Conflict (collision):**
  - Multiple memory locations mapped to the same cache location
  - Solution 1: increase cache size
  - Solution 2: increase associativity

- **Coherence (Invalidation):** other process (e.g., I/O) updates memory
Translation Look-Aside Buffer

- Record recent Virtual Page # to Physical Frame # translation
- If present, have the physical address without reading any page tables !!!
  - Even if the translation involved multiple levels
  - Caches the end-to-end result
- Was invented by Sir Maurice Wilkes – prior to caches
  - When you come up with a new concept, you get to name it!
  - People realized “if it’s good for page tables, why not the rest of the data in memory?”
- On a TLB miss, the page tables may be cached, so only go to memory when both miss

Caching Applied to Address Translation

- Question is one of page locality: does it exist?
  - Instruction accesses spend a lot of time on same page (accesses are sequential)
  - Stack accesses have definite locality of reference
  - Data accesses have less page locality, but still some…
- Can we have a TLB hierarchy?
  - Sure: multiple levels at different sizes/speeds

Physically-Indexed vs Virtually-Indexed Caches

- Physically-Indexed, Physically-Tagged
  - Address handed to cache after translation
  - Page Table in physical memory (so that it can be cached)
  - Benefits:
    » Every piece of data has single place in cache
    » Cache can stay unchanged on context switch
  - Challenges:
    » TLB is in critical path of lookup!
    » Pretty Common today (e.g. x86 processors)
- Virtually-indexed, Virtually-Tagged or Physically-Tagged
  - Address handed to cache before translation
  - Page Table in virtual memory (so that it can be cached); Only last level of Page Table points to physical memory.
  - Benefits:
    » TLB not in critical path of lookup, so system can be faster
  - Challenges:
    » Same data could be mapped in multiple places of cache
    » May need to flush cache on context switch
- We will stick with Physically Indexed Caches for now!

What TLB Organization Makes Sense?

- For Physically Indexed/Tagged, Needs to be really fast
  - Critical path of memory access
    » In simplest view: before the cache
    » Thus, this adds to access time (reducing cache speed)
  - Seems to argue for Direct Mapped or Low Associativity
- However, needs to have very few conflicts!
  - With TLB, the MissTime extremely high! (Page Table traversal)
  - Cost of Conflict (Miss Time) is high
  - Hit Time – dictated by clock cycle
- Thrashing: continuous conflicts between accesses
  - What if use low order bits of virtual page number as index into TLB?
    » First page of code, data, stack may map to same entry
    » Need 3-way associativity at least?
  - What if use high order bits as index?
    » TLB mostly unused for small programs
TLB organization: include protection

- How big does TLB actually have to be?
  - Usually small: 128-512 entries (larger now)
  - Not very big, can support higher associativity
- Small TLBs usually organized as fully-associative cache
  - Lookup is by Virtual Address
  - Returns Physical Address + other info
- What happens when fully-associative is too slow?
  - Put a small (4-16 entry) direct-mapped cache in front
  - Called a “TLB Slice”
- Example for MIPS R3000:

<table>
<thead>
<tr>
<th>Virtual Address</th>
<th>Physical Address</th>
<th>Dirty</th>
<th>Ref</th>
<th>Valid</th>
<th>Access</th>
<th>ASID</th>
</tr>
</thead>
<tbody>
<tr>
<td>0xFA00</td>
<td>0x0003</td>
<td>Y</td>
<td>N</td>
<td>Y</td>
<td>R/W</td>
<td>34</td>
</tr>
<tr>
<td>0x0040</td>
<td>0x0010</td>
<td>N</td>
<td>N</td>
<td>Y</td>
<td>R</td>
<td>0</td>
</tr>
<tr>
<td>0x0041</td>
<td>0x0011</td>
<td>N</td>
<td>Y</td>
<td>Y</td>
<td>R</td>
<td>0</td>
</tr>
</tbody>
</table>

Making physically-indexed caches fast: Fit into Pipeline!

- Example: MIPS R3000 Pipeline

<table>
<thead>
<tr>
<th>Inst Fetch</th>
<th>Dcd/ Reg</th>
<th>ALU / E.A</th>
<th>Memory</th>
<th>Write Reg</th>
</tr>
</thead>
<tbody>
<tr>
<td>TLB</td>
<td>I-Cache</td>
<td>RF</td>
<td>Operation</td>
<td>TLB</td>
</tr>
<tr>
<td>E.A.</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

64 entry, on-chip, fully associative, software TLB fault handler

Virtual Address Space

- ASID
- V. Page Number
- Offset

0xx User segment (caching based on PT/TLB entry)
100 Kernel physical space, cached
101 Kernel physical space, uncached
11x Kernel virtual space

Allows context switching among 64 user processes without TLB flush

Further reducing translation time for physically-indexed caches

- As described, TLB lookup is in serial with cache lookup
  - Consequently, speed of TLB can impact speed of access to cache
- Machines with TLBs go one step further: overlap TLB lookup with cache access
  - Works because offset available early
  - Offset in virtual address exactly covers the "cache index" and "byte select"
  - Thus can select the cached byte(s) in parallel to perform address translation

Overlap of Cache and TLB access

- Here is how this might work with a 4K cache:
  - What if cache size is increased to 8KB?
    - Overlap not complete
    - Need to do something else. See CS152/252
  - As discussed earlier, Virtual Caches would make this faster
    - Tags in cache are virtual addresses
    - Translation only happens on cache misses
What Actually Happens on a TLB Miss?

- Hardware traversed page tables (x86, many others):
  - On TLB miss, hardware in MMU looks at current page table to fill TLB (may walk multiple levels)
    » If PTE valid, hardware fills TLB and processor never knows
    » If PTE marked as invalid, causes Page Fault, after which kernel decides what to do afterwards

- Software traversed Page tables (like MIPS):
  - On TLB miss, processor receives TLB fault
  - Kernel traverses page table to find PTE
    » If PTE valid, fills TLB and returns from fault
    » If PTE marked as invalid, internally calls Page Fault handler

- Most chip sets provide hardware traversal
  - Modern operating systems tend to have more TLB faults since they use translation for many things
  - Examples:
    » shared segments
    » user-level portions of an operating system

Consider weird things that can happen

- What if an instruction has side-effects?
  - Options:
    » Unwind side-effects (easy to restart)
    » Finish off side-effects (messy!)
  - Example 1: mov (sp)+,10
    » What if page fault occurs when write to stack pointer?
    » Did sp get incremented before or after the page fault?
  - Example 2: strcpy (r1), (r2)
    » Source and destination overlap: can’t unwind in principle!
    » IBM S/370 and VAX solution: execute twice – once read-only

- What about “RISC” processors?
  - For instance delayed branches?
    » Example:  bne somewhere
    » id r1,(sp)
    » Restart after page fault: need two PCs, PC and nPC!
  - Delayed exceptions:
    » Example:  div r1, r2, r3  
    » What if takes many cycles to discover divide by zero, but load has already caused page fault?

Precise Exceptions

- Precise $\Rightarrow$ state of the machine is preserved as if program executed up to the offending instruction
  - All previous instructions completed
  - Offending instruction and all following instructions act as if they have not even started
  - Same system code will work on different implementations
  - Difficult in the presence of pipelining, out-of-order execution, ...
    » x86 takes this position

- Imprecise $\Rightarrow$ system software has to figure out what is where and put it all back together
  - Performance goals often lead designers to forsake precise interrupts
    » system software developers, user, markets etc. usually wish they had not done this
  - Modern techniques for out-of-order execution and branch prediction help implement precise interrupts
Current Example: Memory Hierarchy

- Caches (all 64 B line size)
  - L1 I-Cache: 32 KiB/core, 8-way set assoc.
  - L1 D Cache: 32 KiB/core, 8-way set assoc., 4-5 cycles load-to-use, Write-back policy
  - L2 Cache: 1 MiB/core, 16-way set assoc., Inclusive, Write-back policy, 14 cycles latency
  - L3 Cache: 1.375 MiB/core, 11-way set assoc., shared across cores, Non-inclusive victim cache, Write-back policy, 50-70 cycles latency

- TLB
  - L1 ITLB, 128 entries; 8-way set assoc. for 4 KB pages
    » 8 entries per thread; fully associative, for 2 MiB / 4 MiB page
  - L1 DTLB 64 entries; 4-way set associative for 4 KB pages
    » 32 entries; 4-way set associative, 2 MiB / 4 MiB page translations:
      » 4 entries; 4-way associative, 1G page translations:
  - L2 STLB: 1536 entries; 12-way set assoc. 4 KiB + 2 MiB pages
    » 16 entries; 4-way set associative, 1 GiB page translations:

What happens on a Context Switch?

- Need to do something, since TLBs map virtual addresses to physical addresses
  - Address Space just changed, so TLB entries no longer valid!

- Options?
  - Invalidate (“Flush”) TLB: simple but might be expensive
    » What if switching frequently between processes?
  - Include ProcessID in TLB
    » This is an architectural solution: needs hardware

- What if translation tables change?
  - For example, to move page from memory to disk or vice versa...
  - Must invalidate TLB entry!
    » Otherwise, might think that page is still in memory!
  - Called “TLB Consistency”

- Aside: with Virtually-Indexed, Virtually-Tagged cache, need to flush cache!
  - Everyone has their own version of the address “0” and can’t distinguish them
  - This is one advantage of Virtually-Indexed, Physically-Tagged caches..
Page Fault Handling

- The Virtual-to-Physical Translation fails
  - PTE marked invalid, Privilege Level Violation, Access violation, or does not exist
  - Causes an Fault / Trap
    - Not an interrupt because synchronous to instruction execution
  - May occur on instruction fetch or data access
  - Protection violations typically terminate the process
- Other Page Faults engage operating system to fix the situation and retry the instruction
  - Allocate an additional stack page, or
  - Make the page accessible – (Copy on Write),
  - Bring page in from secondary storage to memory – demand paging
- Fundamental inversion of the hardware / software boundary
  - Need to execute software to allow hardware to proceed!

Demand Paging

- Modern programs require a lot of physical memory
  - Memory per system growing faster than 25%-30%/year
- But they don’t use all their memory all of the time
  - 90-10 rule: programs spend 90% of their time in 10% of their code
  - Wasteful to require all of user’s code to be in memory
- Solution: use main memory as “cache” for disk
**Summary (1/2)**

- **The Principle of Locality:**
  - Program likely to access a relatively small portion of the address space at any instant of time.
  - Temporal Locality: Locality in Time
  - Spatial Locality: Locality in Space

- **Three (+1) Major Categories of Cache Misses:**
  - Compulsory Misses: sad facts of life. Example: cold start misses.
  - Conflict Misses: increase cache size and/or associativity
  - Capacity Misses: increase cache size
  - Coherence Misses: Caused by external processors or I/O devices

- **Cache Organizations:**
  - Direct Mapped: single block per set
  - Set associative: more than one block per set
  - Fully associative: all entries equivalent

**Summary (2/2)**

- "Translation Lookaside Buffer" (TLB)
  - Small number of PTEs and optional process IDs (< 512)
  - Often Fully Associative (Since conflict misses expensive)
  - On TLB miss, page table must be traversed and if located PTE is invalid, cause Page Fault
  - On change in page table, TLB entries must be invalidated

- **Demand Paging:** Treating the DRAM as a cache on disk
  - Page table tracks which pages are in memory
  - Any attempt to access a page that is not in memory generates a page fault, which causes OS to bring missing page into memory