

Midterm II
 December 4th, 2006
 CS162: Operating Systems and Systems Programming

Your Name:	
SID Number:	
Circle the letters of CS162 Login	First: a b c d e f g h I j k l m n o p q r s t u v w x y z Second: a b c d e f g h I j k l m n o p q r s t u v w x y z
Discussion Section:	

General Information:

This is a **closed book** exam. You are allowed 1 page of **hand-written** notes (both sides). You have 3 hours to complete as much of the exam as possible. Make sure to read all of the questions first, as some of the questions are substantially more time consuming.

Write all of your answers directly on this paper. *Make your answers as concise as possible.* On programming questions, we will be looking for performance as well as correctness, so think through your answers carefully. If there is something about the questions that you believe is open to interpretation, please ask us about it!

Problem	Possible	Score
1	20	
2	20	
3	27	
4	10	
5	23	
Total		

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3.141592653589793238462643383279502884197169399375105820974944

Problem 1: True/False

In the following, it is important that you *EXPLAIN* your answer in TWO SENTENCES OR LESS (Answers longer than this may not get credit!). Also, answers without an explanation *GET NO CREDIT*.

Problem 1a[2pts]: Memory Mapped I/O devices are accessed with special instructions not used for normal computation.

True / False

Explain:

Problem 1b[2pts]: A single file server can be constructed with a Byzantine Agreement algorithm such that clients are protected against break-ins to the server.

True / False

Explain:

Problem 1c[2pts]: When initiating a secure communication over the web (such as for a bank), you need to trust the public key of a certificate authority that is compiled into your browser.

True / False

Explain:

Problem 1d[2pts]: The limitation in TCP for 65536 ports is due to hardware limitations.

True / False

Explain:

Problem 1e[2pts]: Randomness is essential to achieving good performance from the Ethernet communication algorithm (CSMA/CD).

True / False

Explain:

Problem 1f[2pts]: “Marshalling” is the process by which Byzantine Generals are forced into making good decisions.

True / False

Explain:

Problem 1g[2pts]: RAID-5 is 5 times more reliable than RAID-1.

True / False

Explain:

Problem 1h[2pts]: One advantage of a software TLB is that the same hardware platform can support both forward and inverse page tables.

True / False

Explain:

Problem 1i[2pts]: The *Use* bit in the PTE is required for implementing the Second Chance List algorithm.

True / False

Explain:

Problem 1j[2pts]: The difference between a computer Worm and a computer Virus is that a Virus requires human action to spread.

True / False

Explain:

Problem 2: Virtual Memory, Paging, and Disks

Problem 2a[2pts]: Suppose that we have a 64-bit virtual address split as follows:

9 Bits [Table ID]	13 Bits [Table ID]	13 Bits [Table ID]	13 Bits [Page ID]	16 Bits [Offset]
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Show the format of a PTE complete with bits required to support the clock algorithm, sharing of dynamic libraries, and copy-on-write optimizations.

Problem 2b[3pts]: Explain how to implement the clock algorithm (*not the Second Chance List algorithm*) if the hardware does not directly support *use* bits or *dirty* bits.

Problem 2c[4pts]: For the following problem, assume a hypothetical machine with 4 pages of physical memory and 7 pages of virtual memory. Given the access pattern:

A B C D E A A E C F F G A C G D C F

Indicate in the following table which pages are mapped to which physical pages for each of the following policies. Assume that a blank box matches the element to the left. We have given the FIFO policy as an example.

Access→	A	B	C	D	E	A	A	E	C	F	F	G	A	C	G	D	C	F
FIFO	1	A			E									C				
	2		B			A										D		
	3			C						F								
	4				D							G						
MIN	1																	
	2																	
	3																	
	4																	
LRU	1																	
	2																	
	3																	
	4																	

Problem 2d[4pts]: Suppose that we have a disk with the following parameters:

- 750GB in size
- 12000 RPM, Data transfer rate of 40 Mbytes/s (40×10^6 bytes/sec)
- Average seek time of 8ms
- ATA Controller with 2ms controller initiation time
- A block size of 4Kbytes (4096 bytes)

What is the average time to read a random block from the disk (assuming no queuing at the controller). Show your work. *Hint: there are 4 terms here.*

Problem 2e[2pts]: Given the same parameters from above, assume that the operating system has exploited locality by grouping related blocks together in the filesystem. As a result, the typical access pattern is not as random as in 2d. It typically retrieves 10 blocks sequentially at a time and spends only 1 ms for each seek. What is the average time to read a *single* block now? Show your work.

Problem 2f[5pts]: Assume that there is a queue in front of the controller. Assume that the operating system behaves as in (2e), and that the distribution of service times has $C = 1.5$ (i.e. not quite memoryless). Also assume that requests for blocks arrive via an exponential (memoryless) process with an average arrival rate of λ blocks/second. What is the arrival rate λ such that the queue has an average length of 10 blocks? You don't need to actually compute this number as long as you given an explicit equation for λ . *Hint: use Little's law: $L_q = \lambda T_q$ and solve for utilization.*

Also, remember from basic algebra that $Ax^2 + Bx + C = 0 \Rightarrow x = \frac{-B \pm \sqrt{B^2 - 4AC}}{2A}$.

Problem 3: File Systems

Please keep your answers short (one or two sentences per question-mark). *We may not give credit for long answers.*

Problem 3a[3pts]: Rather than writing updated files to disk immediately when they are closed, many UNIX systems use a delayed *write-behind policy* in which dirty disk blocks are flushed to disk once every 30 seconds. List two advantages and one disadvantage of such a scheme:

Advantage 1:

Advantage 2:

Disadvantage:

Problem 3b[2pts]: Describe a technique that can be used to mitigate the disadvantage of (3a) without losing the advantages of (3a). Be explicit here.

Problem 3c[4pts]: In the following, compare the performance of a FAT file system to a more primitive system that uses “linked allocation” (where file blocks are linked together via pointers in each block) with respect to number of disk accesses. *Be sure to state any assumptions that you are making in terms of what (if anything) is memory resident.*

a) Which is faster for *random* access? Explain. (Assume the file in question is large):

b) Which is faster for *sequential* access? Explain. (Again assume a large file):

Problem 3h[3pts]: The Andrew File System (AFS) solves the above problem (3g) using state information it maintains at the server. What state is kept? How is that state used to solve the problem?

Problem 3i[4pts]: At time XZ, the request queue for a disk contains the following requests in [track:sector] form:

[10:5], [22:9], [20:21], [21:9], [2:10], [40:45], [6:7], [38:9] (in this order).

Assume that the disk head is currently positioned over cylinder 20. What is the sequence of reads under the following head scheduling algorithms?

a) Shortest Seek Time First:

b) SCAN (initially moving upwards):

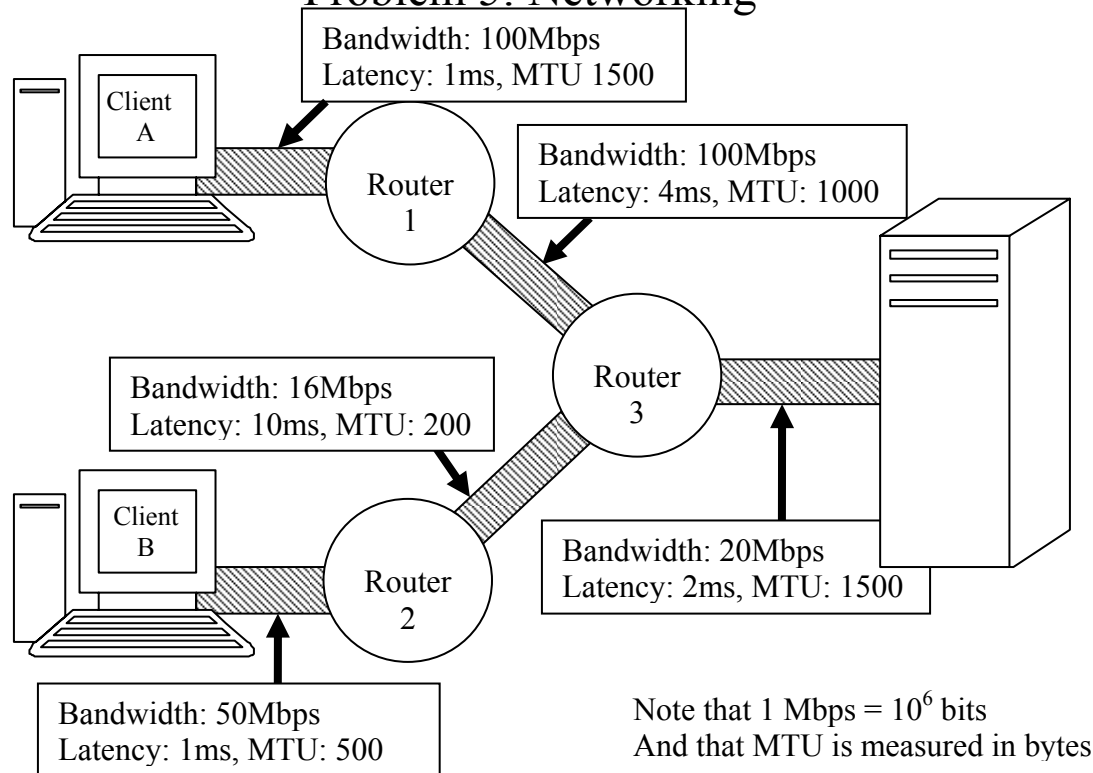
Problem 4: Security

Problem 4a[3pts]: Assume Alice and Bob have never met. Explain how Alice can use a public key infrastructure (PKI) to prove her identity to Bob. *Hint: make sure to prevent replay attacks.*

Problem 4b[3pts]: Explain how to utilize a PKI to establish a private session key between two parties for fast symmetric encryption.

Problem 4c[4pts]: What are two desirable properties for secure hash functions (ignoring the property where half the bits change for small changes in input). Why are these properties important for signatures?

Problem 5: Networking



The above figure illustrates a network in which two clients (Client A and Client B) route packets through the network to the server and to each other. Each link is characterized by its Bandwidth, one-way Latency, and Maximum Transfer Unit (MTU) in *bytes*. All links are full-duplex (can handle traffic in both directions at full bandwidth).

To send data packets, a user-level process on one of the Clients issue a system call. As a result, the OS copies the data first to a kernel buffer, then uses DMA to copy the data to the network controller board. Immediately after receiving all of the data, the network controller generates one or more packets and sends them. A router fully receives a packet, after which it takes 1ms to begin forwarding the packet to the next hop. Assume that no packets will be lost unless otherwise noted.

At the destination, the receiving network controller DMA's the received data to memory. When the last bit arrives, it interrupts the CPU, which copies the data to user space. The CPU then sends back a one-byte acknowledgement to the sending OS. *For simplicity, you may treat acknowledgements as if they are zero bytes long. Assume that a sender uses a window-based protocol unless noted.*

Assume that all DMA operations can overlap with processor actions such as copying and interrupt handling. Also assume that DMA operations do not have to generate interrupts on completion.

System Parameters:

- Time to process an interrupt, $T_{\text{int}} = 100 \mu\text{s}$
- Time to copy or DMA one byte, $T_{\text{copy}} = 100 \text{ ns}$
- Retransmission timeout, $T_{\text{retran}} = 400 \text{ ms}$

Problem 5a[4pts]: Under ideal circumstances (and ignoring interrupt and copying overheads and window size), what is the maximum bandwidth that Client A can send data to the server without causing packets to be dropped (Assuming that the headers are of zero length and that routers can handle data at wire speed)? How about Client B? Explain.

Problem 5b[4pts]: Keeping in mind that TCP/IP involves a total header size of 40 bytes (for TCP + IP), what is the maximum *data* bandwidth that Client A could send to the server through TCP/IP? Explain. What about Client B? Explain. Again, ignore interrupt and copying overheads. *Hint: Assume that a sender runs a protocol that computes the proper maximum packet size to avoid fragmentation within the network.*

Problem 5c[6pts]: Now, assume that copying and interrupt overheads impact the maximum rate at which data can be sent or received at the endpoints (clients or servers). Routers continue to be fully pipelined and senders avoid fragmentation as in 5b. What is the maximum *data* bandwidth that Client A can transfer to the server without dropping packets? Explain. What about Client B? *Hint: make sure to account for the impact of the 40 byte TCP/IP header.*

Problem 5d[5pts]: Now, accounting for copying and interrupt overheads, what is the total time to *reliably* send a single maximal-sized packet from Client A to a user-level process on the server without fragmentation? Account for all latencies (including the reception of an ACK, which can be assumed to be zero length but which will generate an interrupt). Explain your work.

Problem 5e[4pts]: Assume the conditions from 5d. Client A sends a continuous stream of packets to the server (and no other clients are talking to the server). How big should the send window be so that the TCP/IP algorithm will achieve maximum bandwidth without dropping packets? Explain. *Hint: don't forget to account for the 40 bytes of header.*

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